

Network Layer

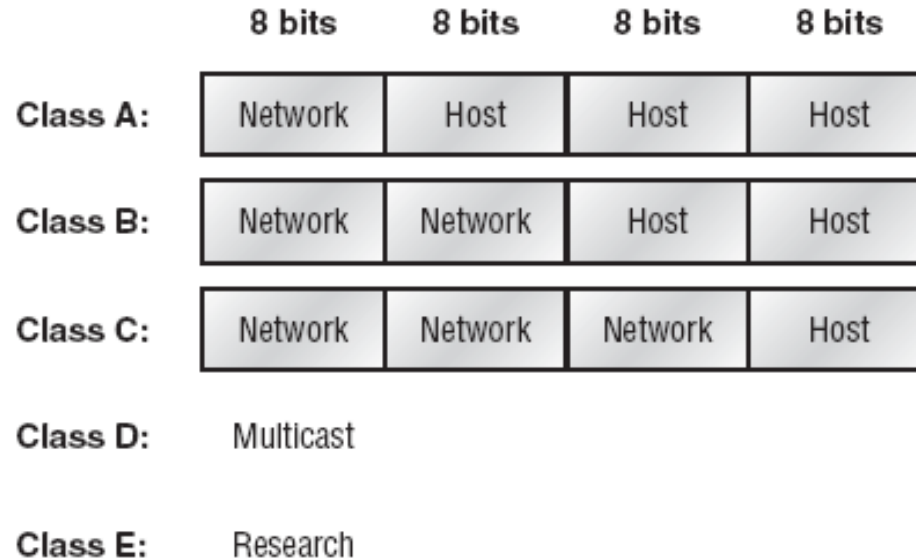
IPv4 addressing

Introduction

- An IP address consists of 32 bits of information.
- These bits are divided into four sections, referred to as *octets* or bytes, each containing 1 byte (8 bits).
- You can depict an IP address using one of three methods:
 - Dotted-decimal, as in 172.16.30.56
 - Binary, as in 10101100.00010000.00011110.00111000
 - Hexadecimal, as in AC.10.1E.38

- The *network address* uniquely identifies each network. Every machine on the same network shares that network address as part of its IP address.
 - For example, 172.16 is the network address.
- The *node address* is assigned to each machine on a network.
 - This part of the address must be unique because it identifies a particular machine
 - This number can also be referred to as a *host address*.
 - IP address 172.16.30.56, the 30.56 is the node address.

Class of IP Address



	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

a. Binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0-127			
Class B	128-191			
Class C	192-223			
Class D	224-239			
Class E	240-255			

b. Dotted-decimal notation

Network Addresses: Special Purpose

Address	Function
Network address of all 0s	Interpreted to mean “this network or segment.”
Network address of all 1s	Interpreted to mean “all networks.”
Network 127.0.0.1	Reserved for loopback tests. Designates the local node and allows that node to send a test packet to itself without generating network traffic.
Node address of all 0s	Interpreted to mean “network address” or any host on specified network.
Node address of all 1s	Interpreted to mean “all nodes” on the specified network; for example, 128.2.255.255 means “all nodes” on network 128.2 (Class B address).
Entire IP address set to all 0s	Used by Cisco routers to designate the default route. Could also mean “any network.”
Entire IP address set to all 1s (same as 255.255.255.255)	Broadcast to all nodes on the current network; sometimes called an “all 1s broadcast” or limited broadcast.

Private IP Addresses

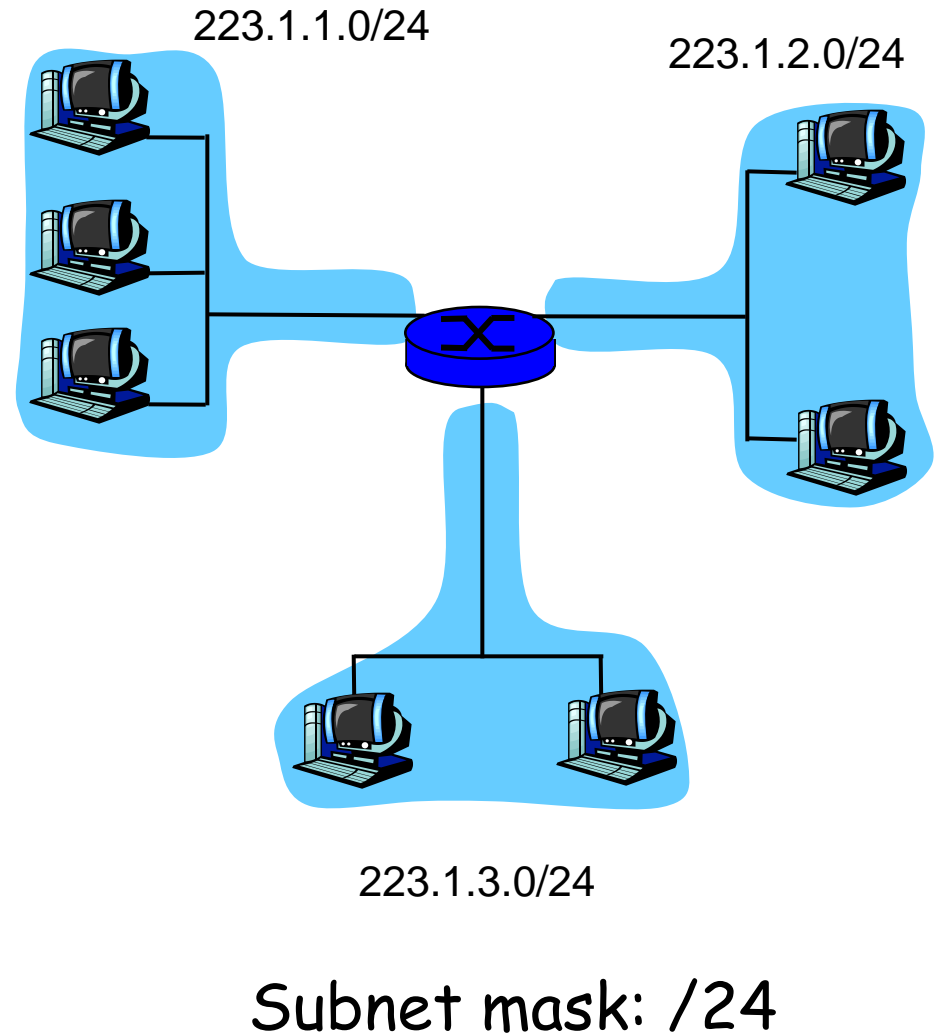
- Private IP addresses
 - used on a private network, but they're not routable through the Internet.
 - Below table shows the reserve private IP address

Address Class	Reserved address space
Class A	10.0.0.0 through 10.255.255.255
Class B	172.16.0.0 through 172.31.255.255
Class C	192.168.0.0 through 192.168.255.255

Subnets

Recipe

- To determine the subnets, detach each interface from its host or router, creating islands of isolated networks. Each isolated network is called a **subnet**.



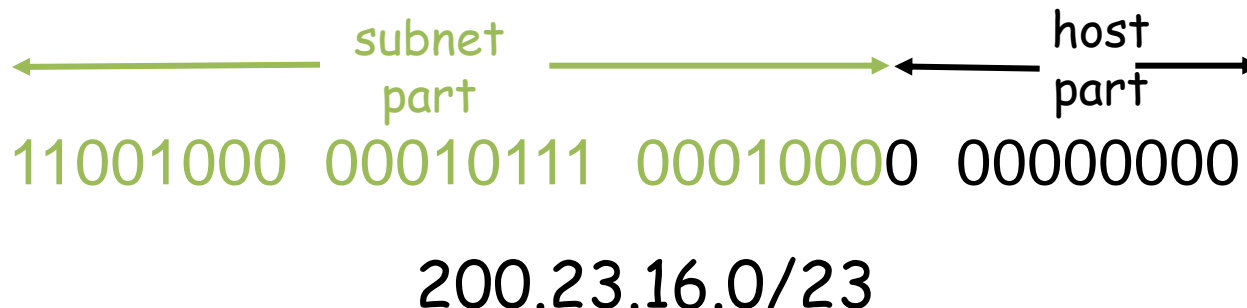
Example

- Class A default subnet mask is 255.0.0.0
 - To a slash notation 255.0.0.0 is considered a /8
- Class B default mask is 255.255.0.0
 - which is a /16
- Class C default mask is 255.255.255.0
 - which is a /24

IP addressing: CIDR

CIDR: Classless InterDomain Routing

- subnet portion of address of arbitrary length
- address format: $a.b.c.d/x$, where x is # bits in subnet portion of address



Subnetting

- What happens if you wanted to take one network address and create six networks from it?
 - From Class B addresses
- **Subnetting** allows you to take one larger network and break it into a bunch of smaller networks

Subnetting Benefit

- Reduced network traffic
- Optimized network performance
- Simplified management
- Facilitated spanning of large geographical distances

Questions for Choosing Subnetting

- How many subnets does the chosen subnet mask produce?
- How many valid hosts per subnet are available?
- What are the valid subnets?
- What's the broadcast address of each subnet?
- What are the valid hosts in each subnet?

Subnets & Hosts

- *How many subnets?*
 - x = number of subnets. x is the number of masked bits, or the 1s. For example, 11000000, the number of 1s gives us 2^2 subnets. There are 4 subnets.
- *How many hosts per subnet?*
 - $2^y - 2$ = number of hosts per subnet. y is the number of
 - unmasked bits, or the 0s. For example, in 11000000, the number of 0s gives us $2^6 - 2$ hosts.
 - You need to subtract 2 for the subnet address and the broadcast address, which are not valid hosts.

Valid Subnets

- *What are the valid subnets?*
 - $256 - \text{subnet mask} = \text{block size, or increment number.}$
 - An example would be $256 - 192 = 64$. The block size of a 192 mask is always 64.
 - Start counting at zero in blocks of 64 until you reach the subnet mask value and these are your subnets. 0, 64, 128, 192.

Broadcast Address

- *What's the broadcast address for each subnet?*
 - The broadcast address is always the number right before the next subnet.
 - For example, the 0 subnet has a broadcast address of 63 because the next subnet is 64. The 64 subnet has a broadcast address of 127 because the next subnet is 128. And so on.

Valid Hosts

- *What are the valid hosts?* Valid hosts are the numbers between the subnets, omitting the all 0s and all 1s.
 - For example, if 64 is the subnet number and 127 is the broadcast address, then 65–126 is the valid host range—it's *always* the numbers between the subnet address and the broadcast address.

Examples: Class C Addresses

- Class C network address 192.168.10.0 with subnet 255.255.255.128 (/25)
 - Since 128 is 10000000 in binary, there is only 1 bit for subnetting and 7 bits for hosts.
 - 192.168.10.0 = Network address
 - 255.255.255.128 = Subnet mask

So....

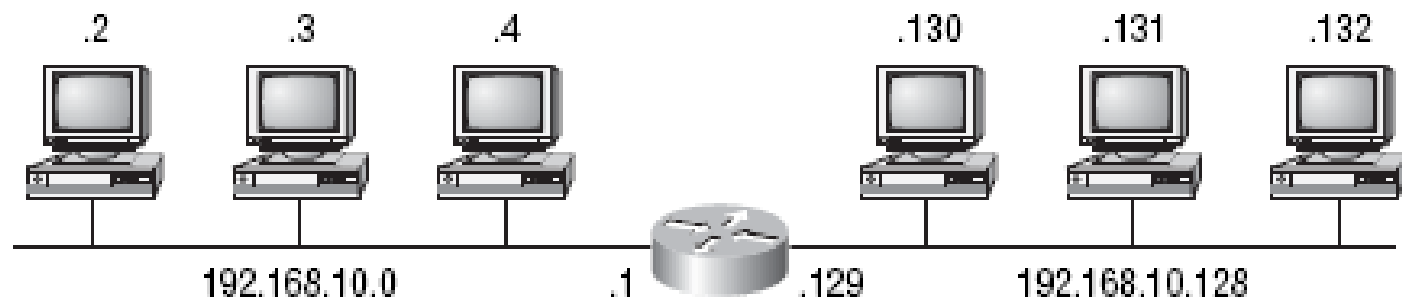
- *How many subnets?*
 - Since 128 is 1 bit on (**1**0000000), the answer would be $2^1 = 2$.
- *How many hosts per subnet?*
 - We have 7 host bits off (**1**0000000), so the equation would be $2^7 - 2 = 126$ hosts.
- *What are the valid subnets?* $256 - 128 = 128$.
 - Start at zero and count in our block size, so our subnets are 0, 128.

So...

- *What's the broadcast address for each subnet?*
 - The number right before the value of the next subnet. For the zero subnet, the next subnet is 128, so the broadcast of the 0 subnet is 127.
- *What are the valid hosts?*

Subnet	0	128
First host	1	129
Last host	126	254
Broadcast	127	255

So... The network



Subnetting Class B Addresses

- Class B address 172.16.0.0 Subnet mask 255.255.192.0 (/18)
 - *Subnets?* $2^2 = 4$.
 - *Hosts?* $2^{14} - 2 = 16,382$ (6 bits in the third octet, and 8 in the fourth).
 - *Valid subnets?* $256 - 192 = 64$. 0, 64, 128, 192.
Remember that the subnetting is performed in the third octet, so the subnet numbers are really 0.0, 64.0, 128.0, and 192.0,

- *Broadcast address for each subnet?*
- *Valid hosts?*

Subnet	0.0	64.0	128.0	192.0
First host	0.1	64.1	128.1	192.1
Last host	63.254		127.254	191.254
	255.254			
Broadcast	63.255		127.255	191.255
	255.255			

IP addresses: how to get one?

Q: How does *host* get IP address?

- hard-coded by system admin in a file
 - Wintel: control-panel->network->configuration->tcp/ip->properties
 - UNIX: /etc/rc.config
- **DHCP: Dynamic Host Configuration Protocol:**
dynamically get address from as server
 - “plug-and-play”

DHCP: Dynamic Host Configuration Protocol

Goal: allow host to *dynamically* obtain its IP address from network server when it joins network

Can renew its lease on address in use

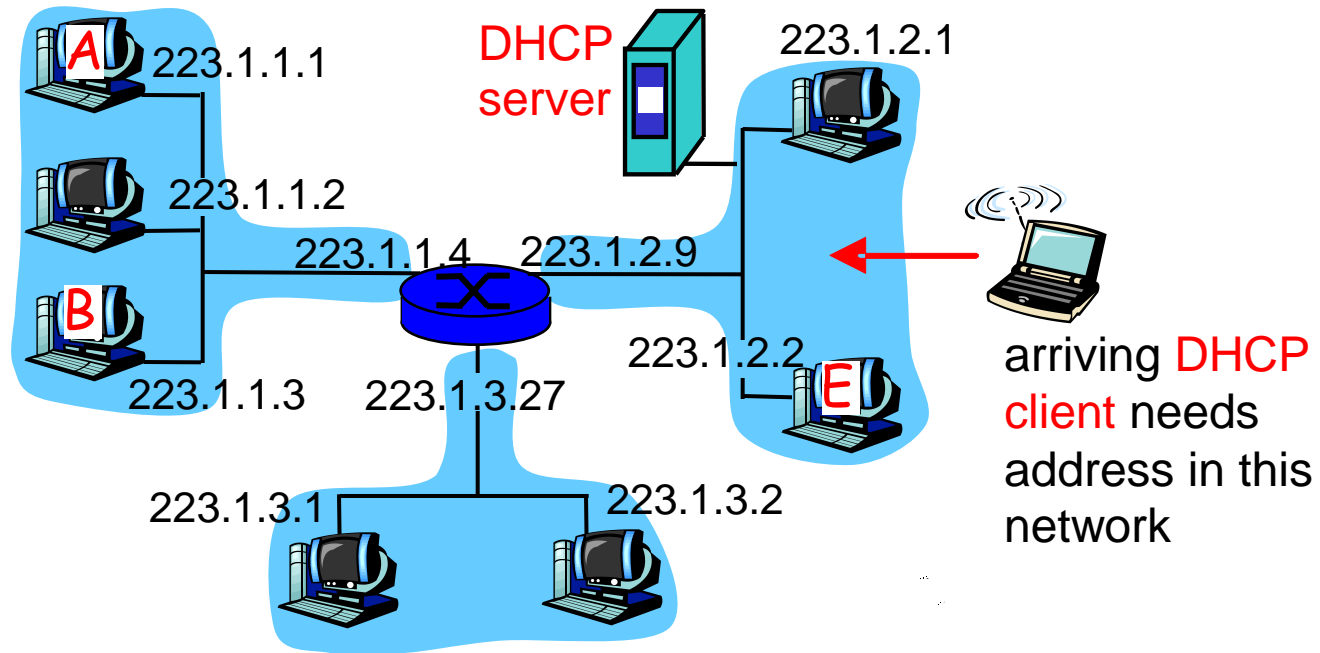
Allows reuse of addresses (only hold address while connected an “on”

Support for mobile users who want to join network (more shortly)

DHCP overview:

- host broadcasts “DHCP discover” msg
- DHCP server responds with “DHCP offer” msg
- host requests IP address: “DHCP request” msg
- DHCP server sends address: “DHCP ack” msg

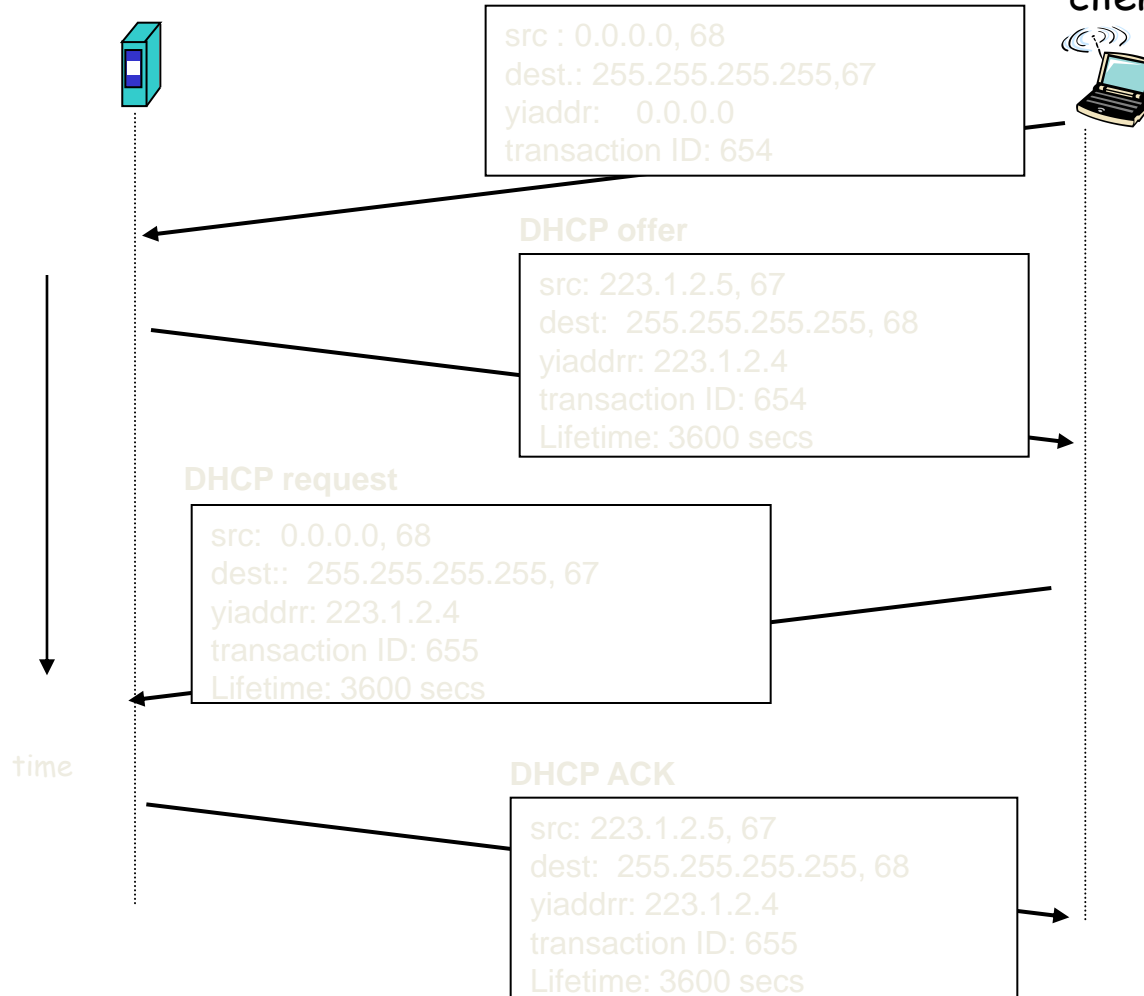
DHCP client-server scenario



DHCP client-server scenario

DHCP server: 223.1.2.5

arriving
client



IP addresses: how to get one?

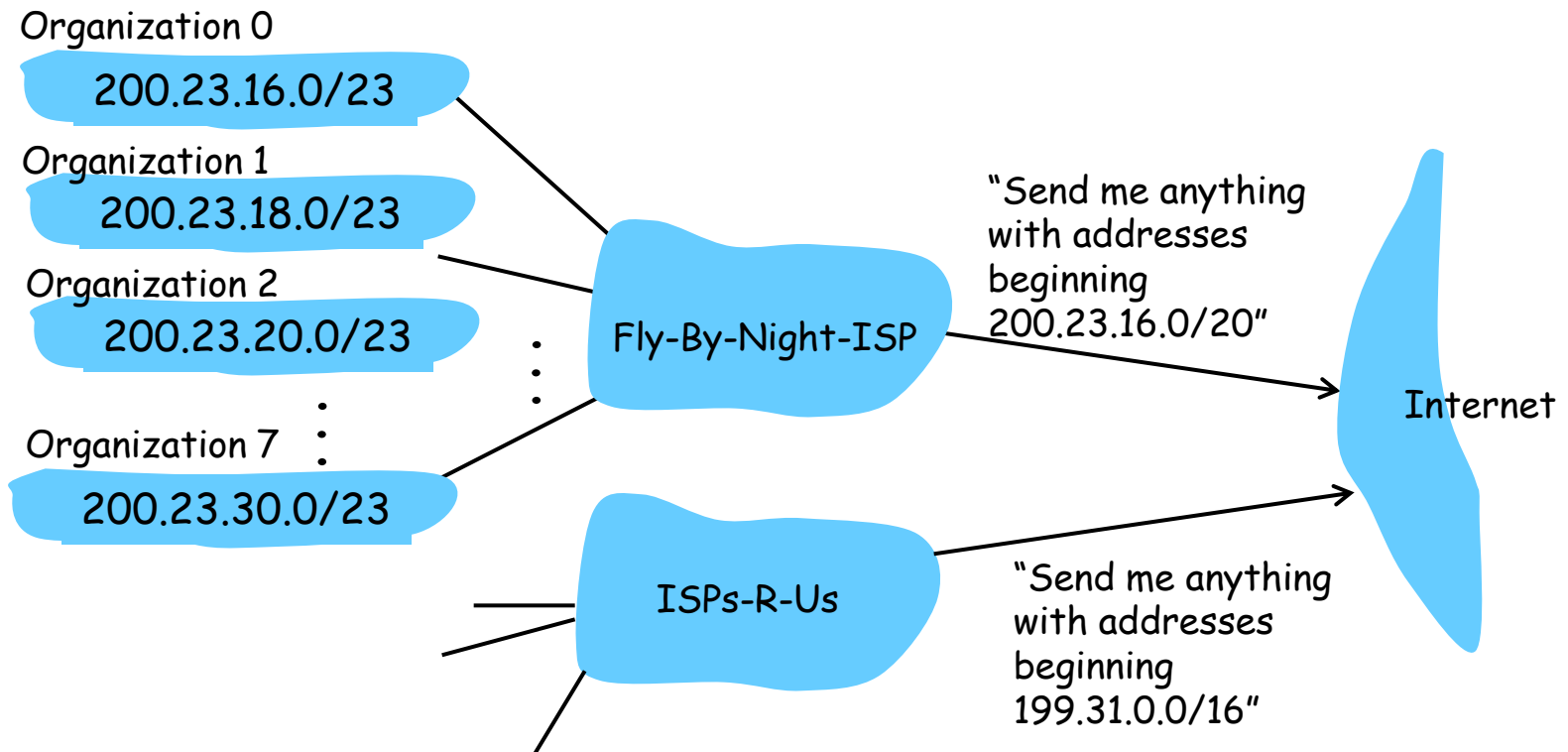
Q: How does *network* get subnet part of IP addr?

A: gets allocated portion of its provider ISP's address space

ISP's block	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/20
Organization 0	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/23
Organization 1	<u>11001000</u>	<u>00010111</u>	<u>00010010</u>	00000000	200.23.18.0/23
Organization 2	<u>11001000</u>	<u>00010111</u>	<u>00010100</u>	00000000	200.23.20.0/23
...
Organization 7	<u>11001000</u>	<u>00010111</u>	<u>00011110</u>	00000000	200.23.30.0/23

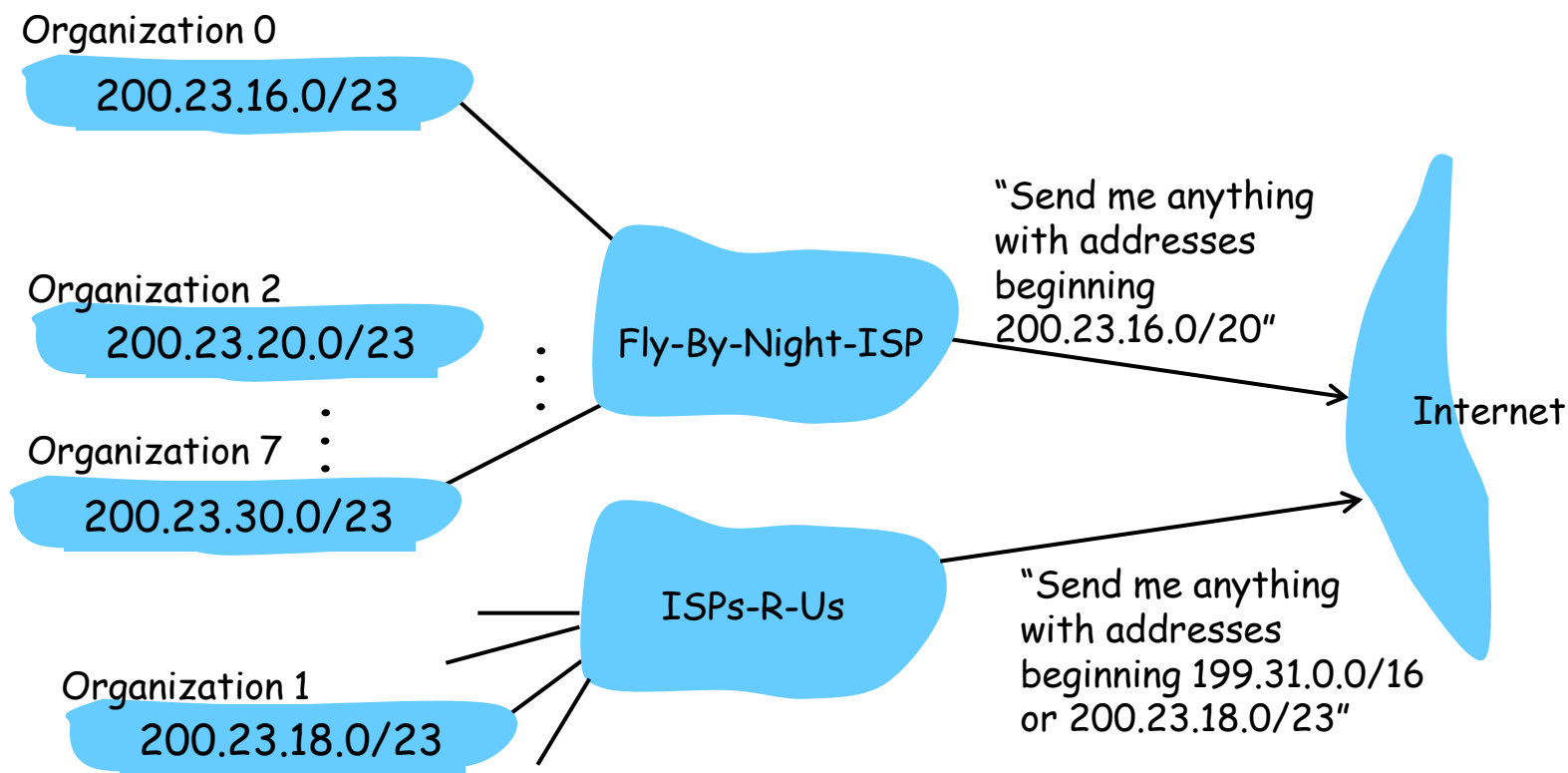
Hierarchical addressing: route aggregation

Hierarchical addressing allows efficient advertisement of routing information:



Hierarchical addressing: more specific routes

ISPs-R-Us has a more specific route to Organization 1



IP addressing: the last word...

Q: How does an ISP get block of addresses?

A: ICANN: Internet Corporation for Assigned
Names and Numbers

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes

Network Layer

ICMP

ICMP: Internet Control Message Protocol

- used by hosts & routers to communicate network-level information
 - error reporting: unreachable host, network, port, protocol
 - echo request/reply (used by ping)
- network-layer “above” IP:
 - ICMP msgs carried in IP datagrams
- **ICMP message:** type, code plus first 8 bytes of IP datagram causing error

<u>Type</u>	<u>Code</u>	<u>description</u>
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

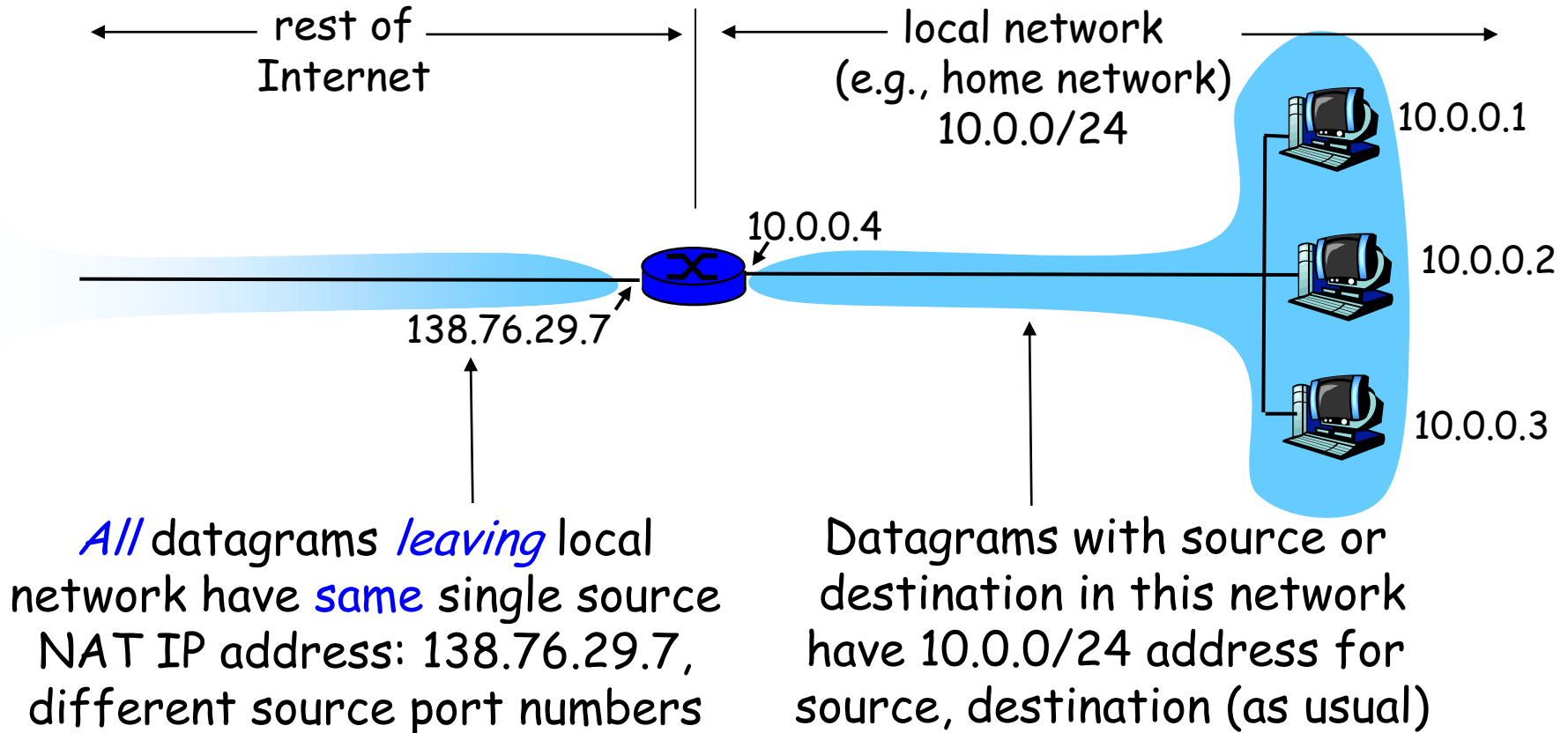
Traceroute and ICMP

- Source sends series of UDP segments to dest
 - First has TTL =1
 - Second has TTL=2, etc.
 - Unlikely port number
 - When nth datagram arrives to nth router:
 - Router discards datagram
 - And sends to source an ICMP message (type 11, code 0)
 - Message includes name of router& IP address
 - When ICMP message arrives, source calculates RTT
 - Traceroute does this 3 times
- Stopping criterion
- UDP segment eventually arrives at destination host
 - Destination returns ICMP “host unreachable” packet (type 3, code 3)
 - When source gets this ICMP, stops.

Network Layer

NAT

NAT: Network Address Translation



NAT: Network Address Translation

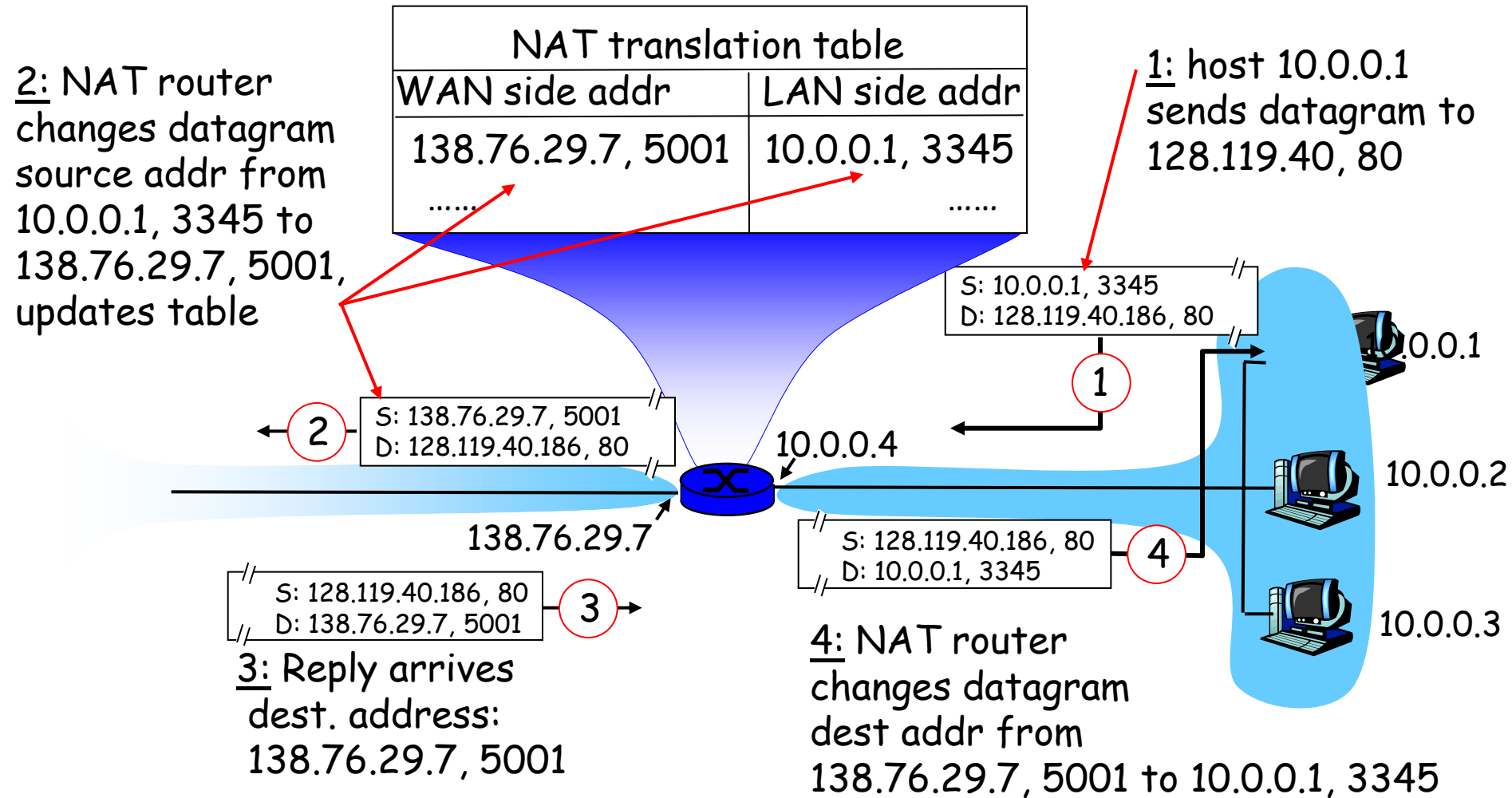
- **Motivation:** local network uses just one IP address as far as outside world is concerned:
 - no need to be allocated range of addresses from ISP:
 - just one IP address is used for all devices
 - can change addresses of devices in local network without notifying outside world
 - can change ISP without changing addresses of devices in local network
 - devices inside local net not explicitly addressable, visible by outside world (a security plus).

NAT: Network Address Translation

Implementation: NAT router must:

- *outgoing datagrams: replace* (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
... remote clients/servers will respond using (NAT IP address, new port #) as destination addr.
- *remember (in NAT translation table)* every (source IP address, port #) to (NAT IP address, new port #) translation pair
- *incoming datagrams: replace* (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

NAT: Network Address Translation



NAT: Network Address Translation

- 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
 - routers should only process up to layer 3
 - violates end-to-end argument
 - NAT possibility must be taken into account by app designers, eg, P2P applications
 - address shortage should instead be solved by IPv6

Network Layer

IPv6

IPv6

- Initial motivation: 32-bit address space soon to be completely allocated.
- Additional motivation:
 - header format helps speed processing/forwarding
 - header changes to facilitate QoS

IPv6 datagram format:

- fixed-length 40 byte header
- no fragmentation allowed

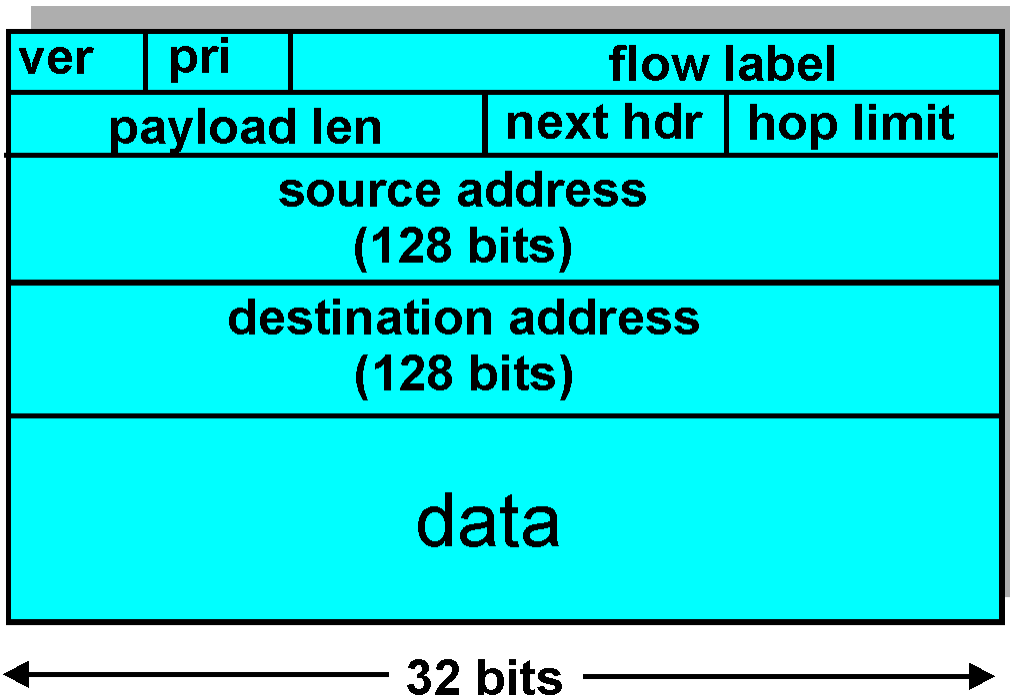
IPv6 Header (Cont)

Priority: identify priority among datagrams in flow

Flow Label: identify datagrams in same "flow."

(concept of "flow" not well defined).

Next header: identify upper layer protocol for data



Other Changes from IPv4

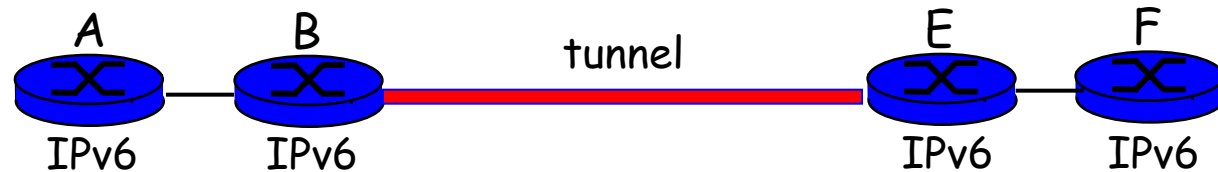
- *Checksum*: removed entirely to reduce processing time at each hop
- *Options*: allowed, but outside of header, indicated by “Next Header” field
- *ICMPv6*: new version of ICMP
 - additional message types, e.g. “Packet Too Big”
 - multicast group management functions

Transition From IPv4 To IPv6

- Not all routers can be upgraded simultaneous
 - no “flag days”
 - How will the network operate with mixed IPv4 and IPv6 routers?
- *Tunneling*: IPv6 carried as payload in IPv4 datagram among IPv4 routers

Tunneling

Logical view:



Physical view:

